

FIG 1

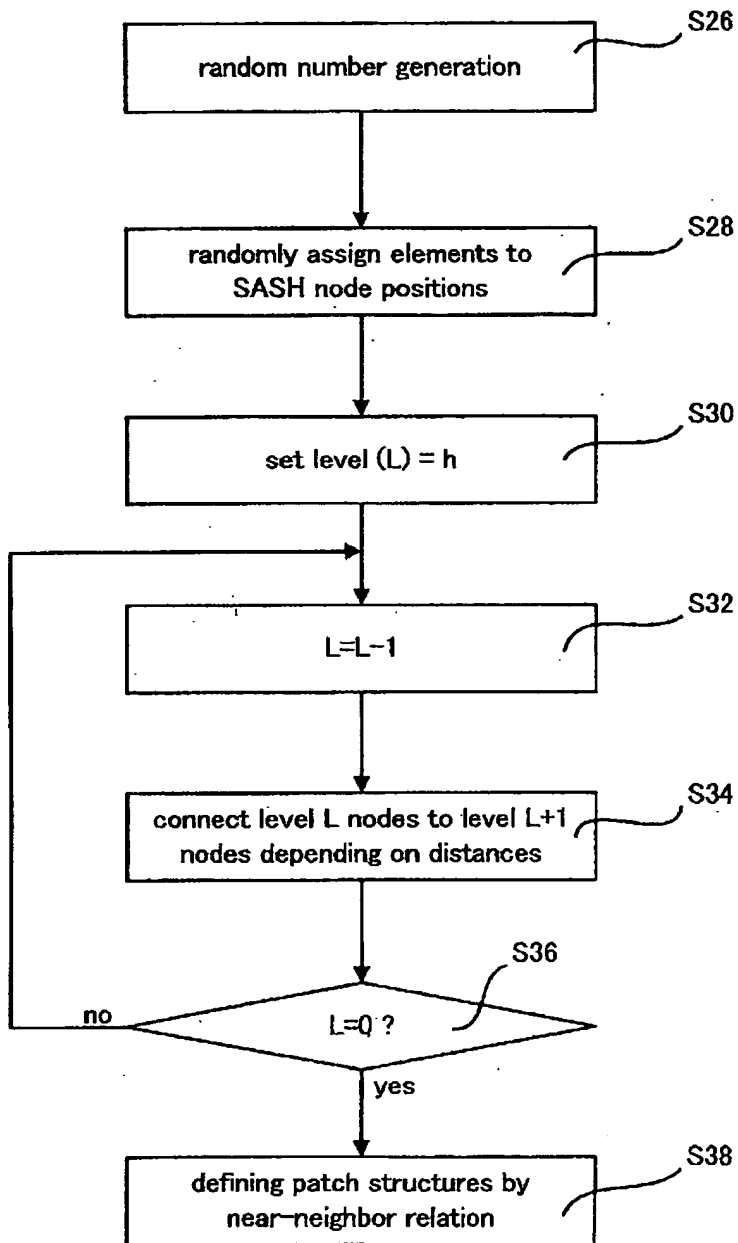


FIG 2

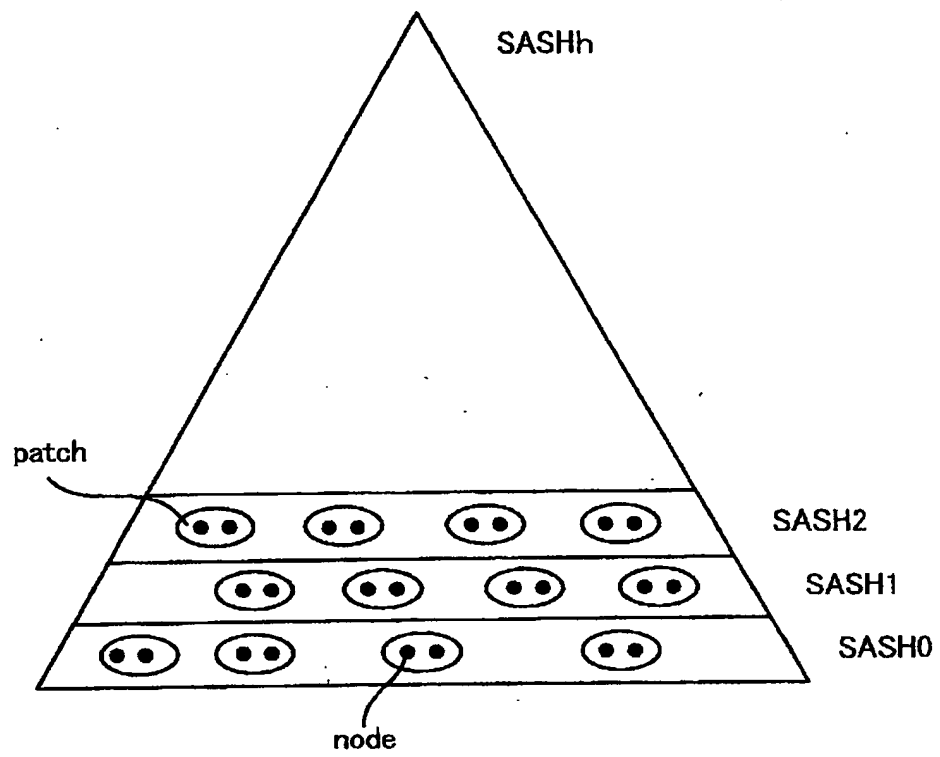


FIG 3

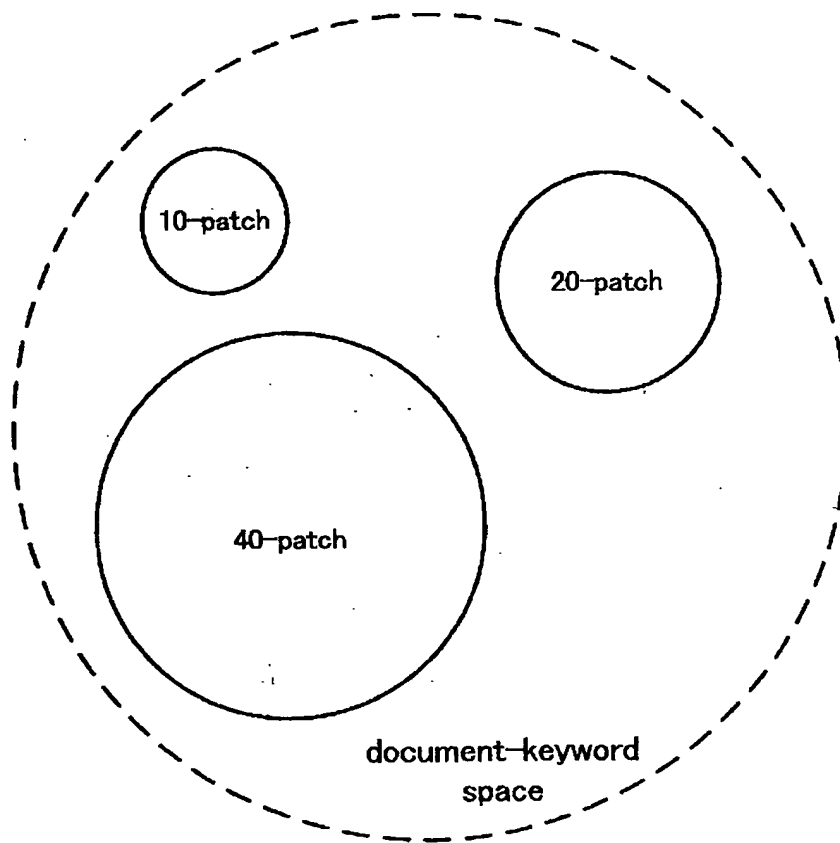
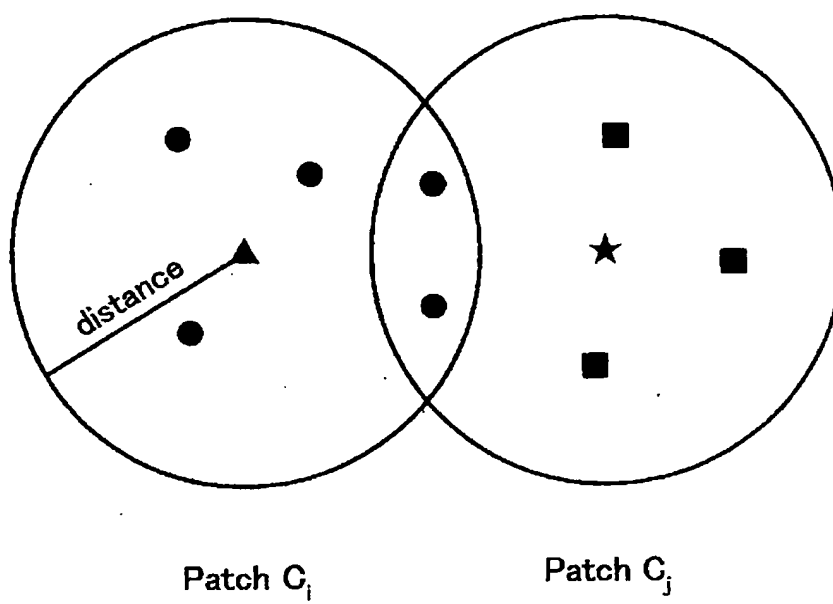


FIG 4



$$\text{CONF}(C_i, C_j) = 2/5 = 40\%$$

FIG 5

Profile (query q ; maximum patch size m): SCONE list SCONFL

{Let QNL be the m -patch precomputed for query q .}

{Let NNL be a list of the m -patches precomputed for every element of QNL .}

{Initially, $w.count = \emptyset$ is assumed for every element v in the data set.}

1. $score \leftarrow 0$;
 {Initially, no query neighbors are in the current patch.}
 for $i = 1$ to m do
2. $QNL[i].count \leftarrow 0$;
 end for
 for $i = 1$ to m do
 {Retrieve the number of times $QNL[i]$ has been encountered as an external neighbor so far.}
3. $score \leftarrow score + QNL[i].count$;
 {Indicate that henceforth $QNL[i]$ is in the current i -patch.}
4. $QNL[i].count \leftarrow present$;
 for $j = 1$ to $i - 1$ do
5. $w \leftarrow NNL[j, i]$;
 if $w.count = present$ then
6. $score \leftarrow score + 1$;
7. else if $w.count \geq 0$ then
 $w.count \leftarrow w.count + 1$;
8. end if
9. $w \leftarrow NNL[i, j]$;
 if $w.count = present$ then
10. $score \leftarrow score + 1$;
11. else if $w.count \geq 0$ then
 $w.count \leftarrow w.count + 1$;
12. end if
13. end for
14. $w \leftarrow NNL[i, i]$;
 if $w.count = present$ then
15. $score \leftarrow score + 1$;
16. else if $w.count \geq 0$ then
 $w.count \leftarrow w.count + 1$;
17. end if
18. $SCONFL[i] = score/i^2$;
19. end for
 {Reset the counts to their default value.}
20. for $i = 1$ to m do
21. $QNL[i].count \leftarrow \emptyset$;
22. end for

FIG 6

SASH #	patch list	CONFL	SCONFL	RSCONFL
0	NN($R_{1,0,0^*m}$),
1	NN($R_{1,1,0^*m}$),
...
h	NN($R_{1,0,0^*m}$),

FIG 7

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RefineProfile (query  $q$ ;
    inner patch size  $k_I$ ;
    outer patch size  $k_O$ ): reordered query  $k_I$ -patch  $RQNL$ 
{Let  $QNL$  be the  $k_O$ -patch precomputed for query  $q$ .}
{Let  $NNL$  be a list of the  $k_O$ -patches precomputed for every element of  $QNL$ .}
{Initially,  $v.inpatch = false$  is assumed for every element  $v$  in the data set.}
{Identify the inner patch members.}
for  $i = 1$  to  $k_I$  do
1.    $QNL[i].inpatch \leftarrow true$ ;
end for
{Initialize the confidence value  $CONF_c$  of every patch element to zero.}
for  $i = 1$  to  $k_O$  do
2.    $CONF_c[i] \leftarrow 0$ ;
end for
{For each element of the outer patch, count the number of elements
of their  $k$ -nearest-neighbor sets shared with that of  $q$ .}
for  $i = 1$  to  $k_O$  do
    for  $j = 1$  to  $k_I$  do
3.        $w \leftarrow NNL[i, j]$ ;
        if  $w.inpatch = true$  then
4.            $CONF_c[i] \leftarrow CONF_c[i] + 1$ ;
        end if
    end for
5.    $CONF_c[i] \leftarrow CONF_c[i] / k_O$ ;
end for
{Reorder the outer patch elements according to their confidence values, from highest to lowest.}
6.  $RQNL \leftarrow sort(QNL, CONF_c, k_O)$ ;
{Reset the patch membership indicators to their default values.}
for  $i = 1$  to  $k_I$  do
7.    $QNL[i].inpatch \leftarrow false$ ;
end for

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FIG 8

PatchCluster (data set S ;

RSCM parameters $a, b, m = \varphi(b)$;

Thresholds $\alpha, \beta, \gamma, \delta$): query cluster graph G

1. Randomly partition the set S into subsets S_t of approximate size $\frac{|S|}{3^t}$, for $0 \leq t \leq h = \lceil \log_3 |S| \rceil$.

2. For all $0 \leq t \leq h$ do:

(a) For every element $v \in S_t$, compute nearest-neighbor patches $NN(R_t, v, m)$, where $R_t = \bigcup_{i \geq t} S_i$.

(b) For each element $v_{t,i} \in S_t$, compute the optimal query cluster size $k(v_{t,i})$ maximizing $RSCONF(NN(R_t, v_{t,i}, k), \varphi)$, for values of k between a and b

The ranked collection of patches

$$C_t = \{C_{t,i} | i < j \Rightarrow RSCONF(C_{t,i}, \varphi) \geq RSCONF(C_{t,j}, \varphi)\}$$

form the candidates for the query clusters associated with sample $R_t \subseteq S$, where $C_{t,i} = NN(R_t, v_{t,i}, k(v_{t,i}))$ and $C_{t,j} = NN(R_t, v_{t,j}, k(v_{t,j}))$.

(c) Let Q_t be a list of patches of C_t that have been confirmed as query clusters of R_t . Initially, Q_t is empty.

(d) For all $1 \leq i \leq |C_t|$ do:

i. If $RSCONF(C_{t,i}, \varphi) < \alpha$, then break from the loop.

ii. For all $w \in C_{t,i}$ do: if $NN(R_t, w, k) \notin |C_t|$ for any value of k , or failing that, if $\max\{CONF(NN(R_t, w, k), C_{t,i}), CONF(C_{t,i}, NN(R_t, w, k))\} < \beta$, then add $C_{t,i}$ to the list Q_t .

3. Let h' be the largest index for which $|Q_{h'}| > 0$. Let $\{C_{t,j}\}$ be the set of patches comprising Q_t , where $C_{t,j} = NN(R_t, q_{t,j}, k(q_{t,j}))$, for all $0 \leq t \leq h'$. Initialize the node set of the query cluster graph G to contain these patches, one patch per node.

4. For all $\delta \leq t \leq h'$, all $1 \leq j \leq |Q_t|$, and all $\max\{0, t - \delta\} \leq s \leq t$, do:

(a) Compute $C'_{t,j} = NN(R_s, q_{t,j}, 2^{t-s}k(q_{t,j}))$.

(b) For all $1 \leq i \leq |Q_s|$, if $C_{s,i} \neq C'_{t,j}$ and $\max\{CONF(C_{s,i}, C'_{t,j}), CONF(C'_{t,j}, C_{s,i})\} \geq \gamma$, then introduce the edges $(C_{s,i}, C_{t,j})$ and $(C_{t,j}, C_{s,i})$ into G , with weights $CONF(C_{s,i}, C'_{t,j})$ and $CONF(C'_{t,j}, C_{s,i})$, respectively.

FIG 9

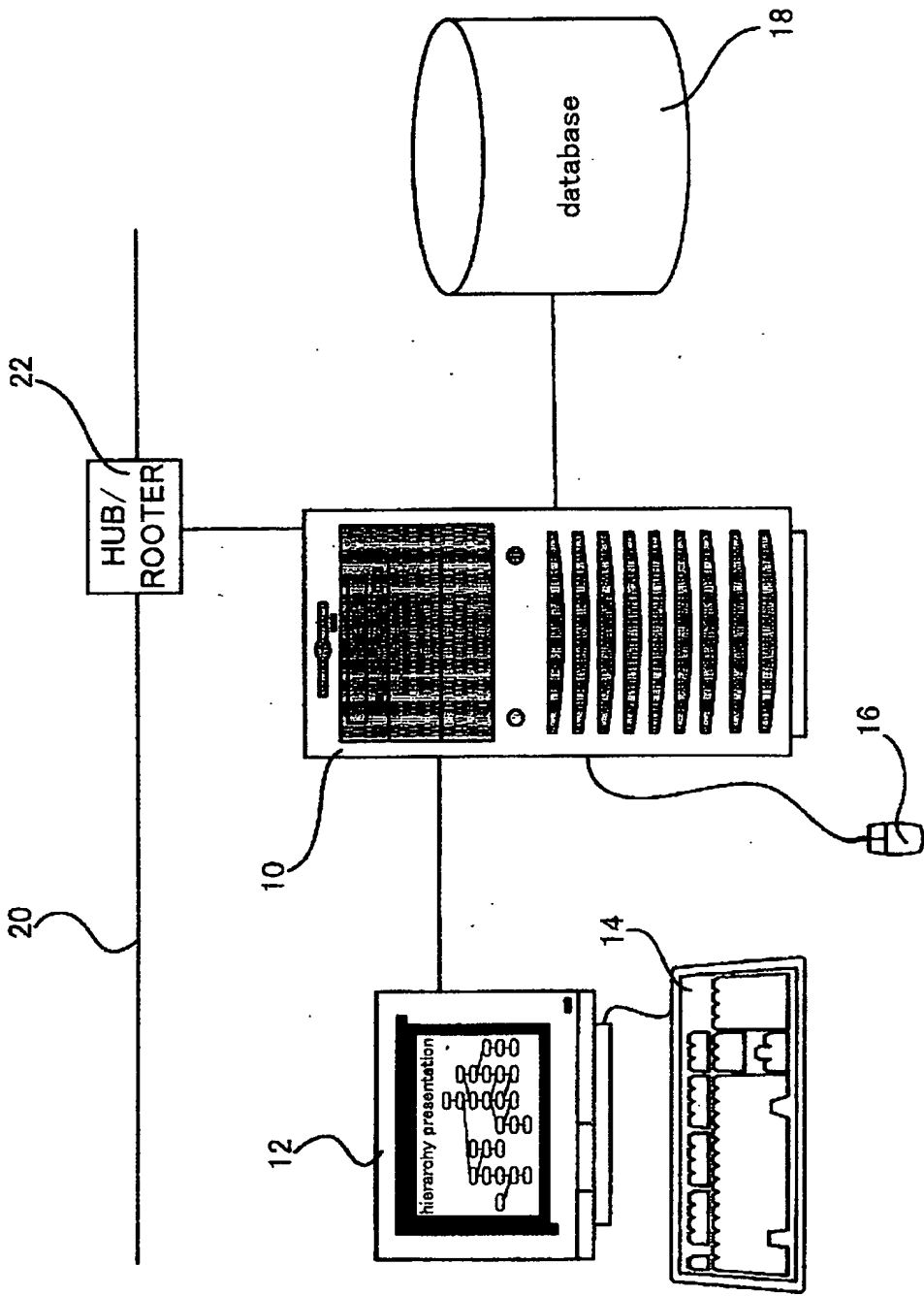


FIG 10

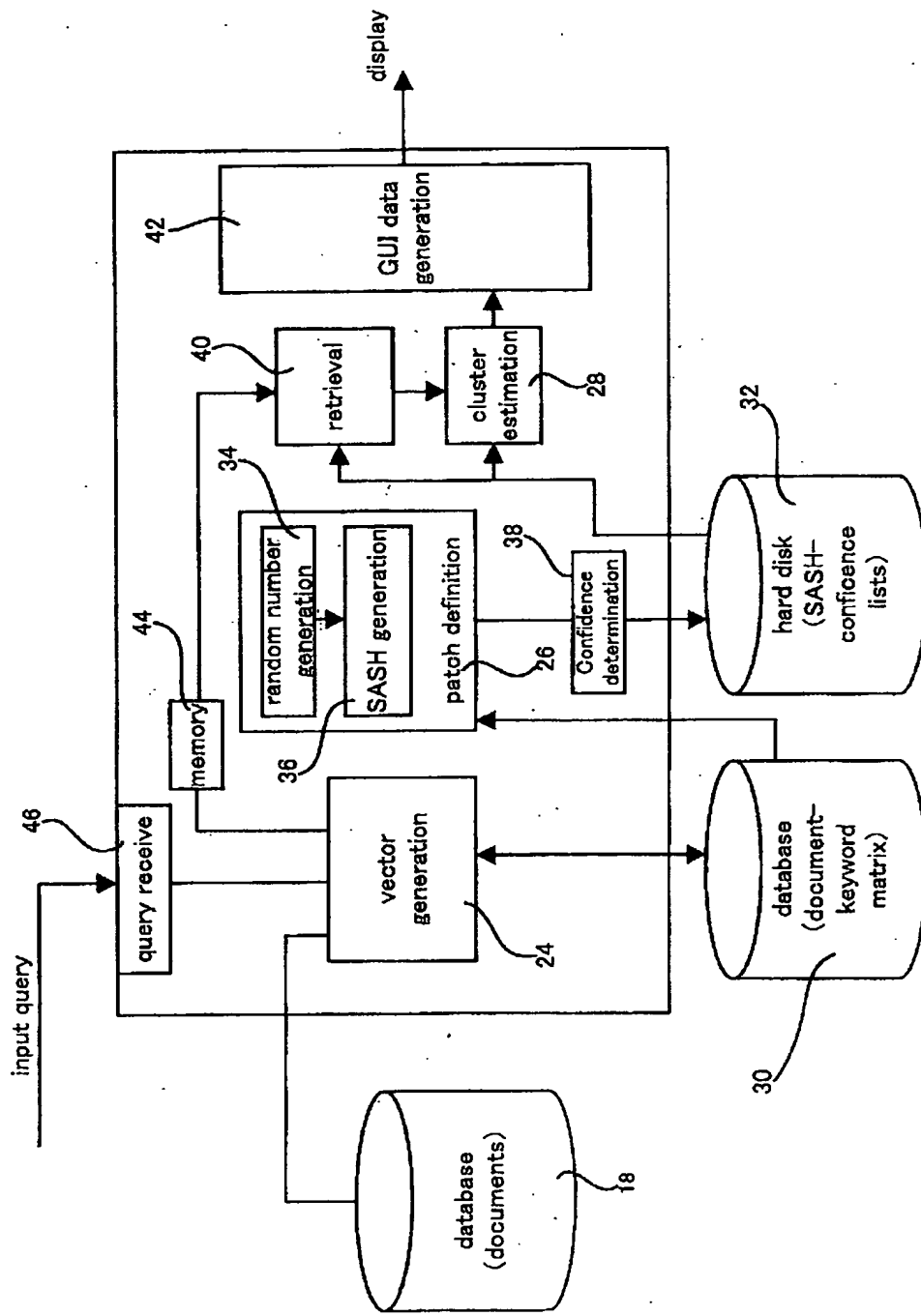


FIG 11

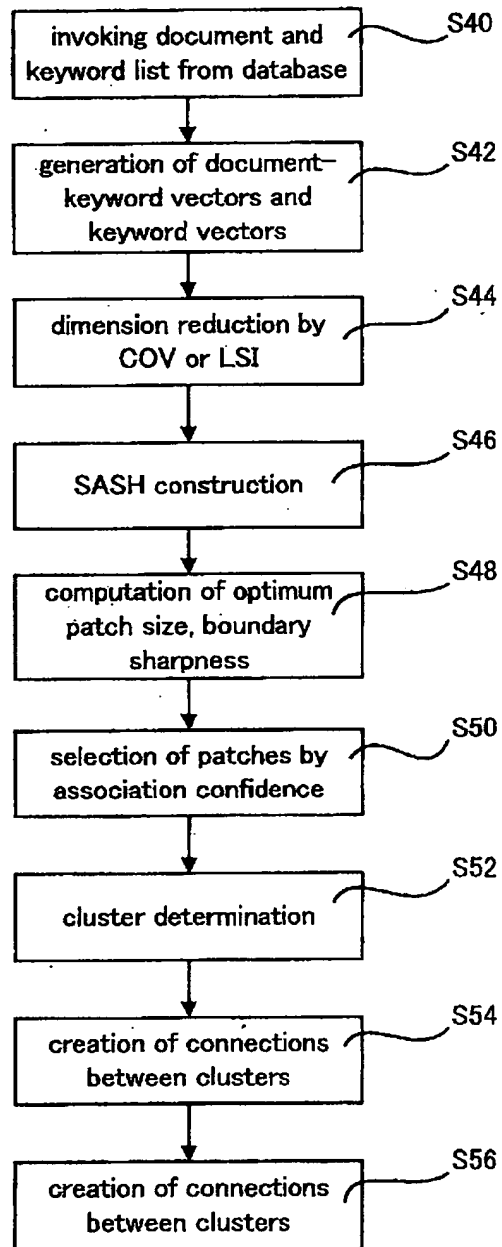


FIG 12

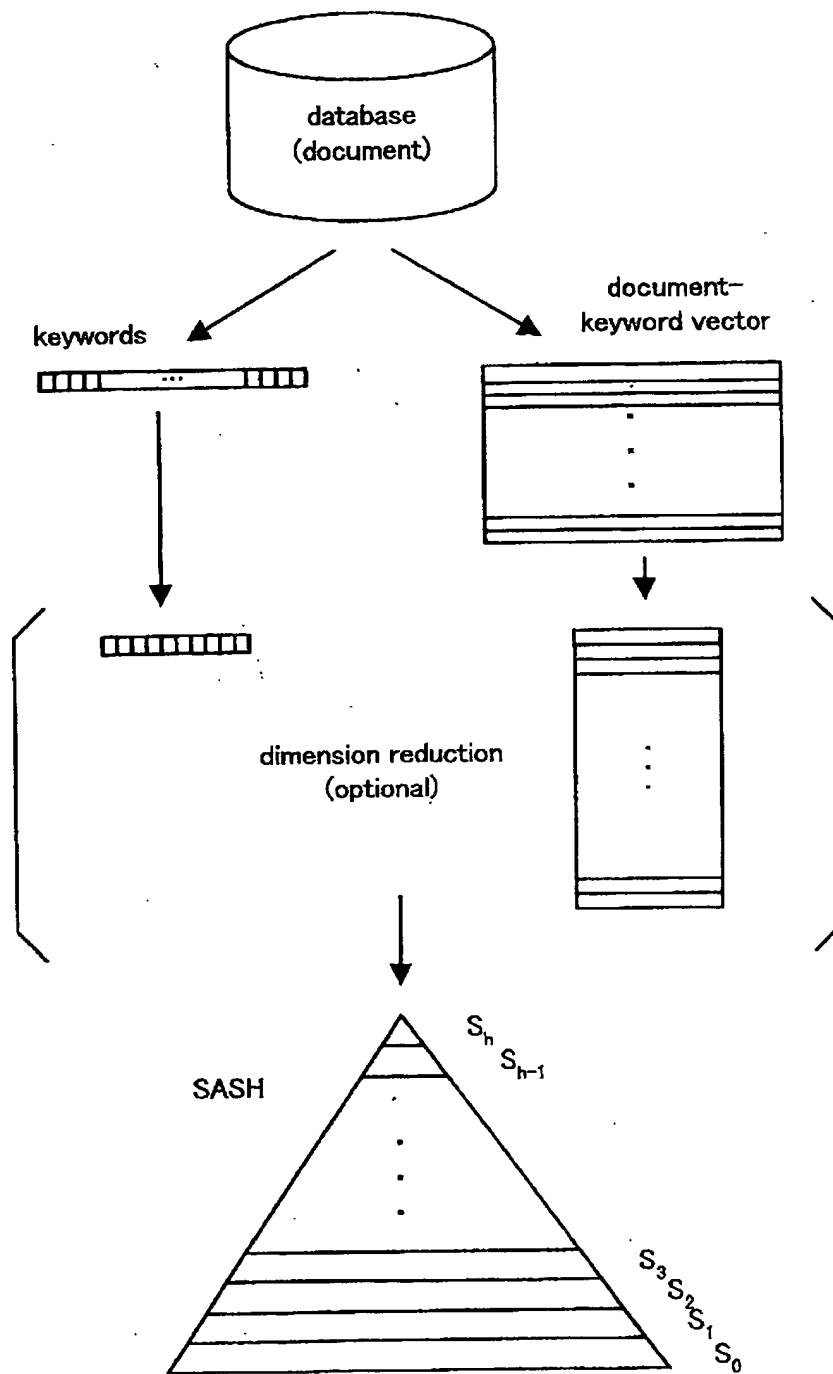
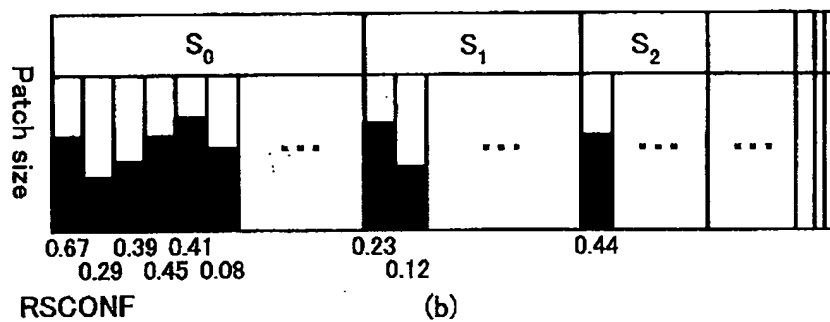


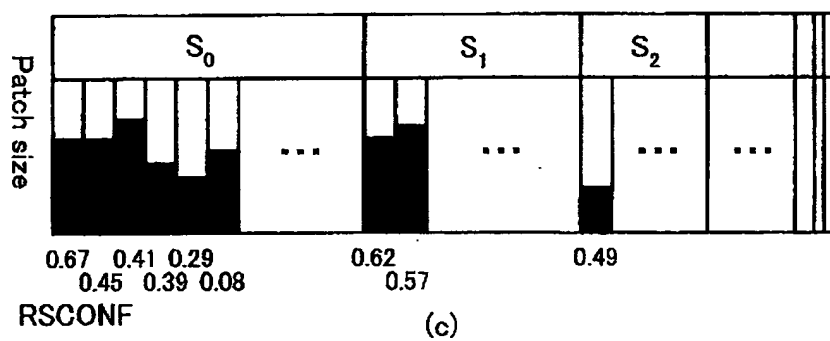
FIG13

S_0						S_1		S_2			
$NN(R_0, V_{0,0}, m)$	$NN(R_0, V_{0,1}, m)$	$NN(R_0, V_{0,2}, m)$	$NN(R_0, V_{0,3}, m)$	$NN(R_0, V_{0,4}, m)$	$NN(R_0, V_{0,5}, m)$	$NN(R_1, V_{0,0}, m)$	$NN(R_1, V_{0,1}, m)$	$NN(R_2, V_{0,0}, m)$	
...						

(a)



(b)



(c)

FIG 14

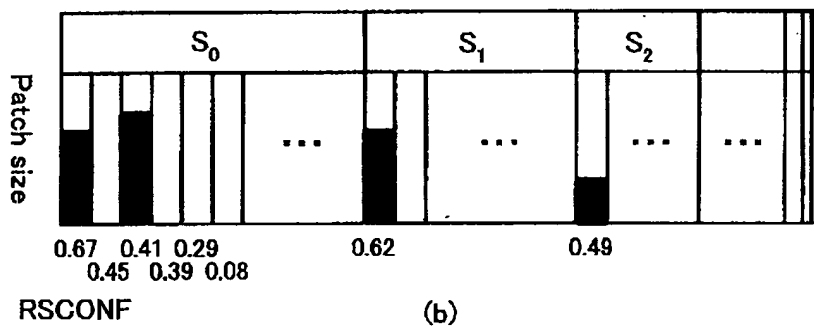
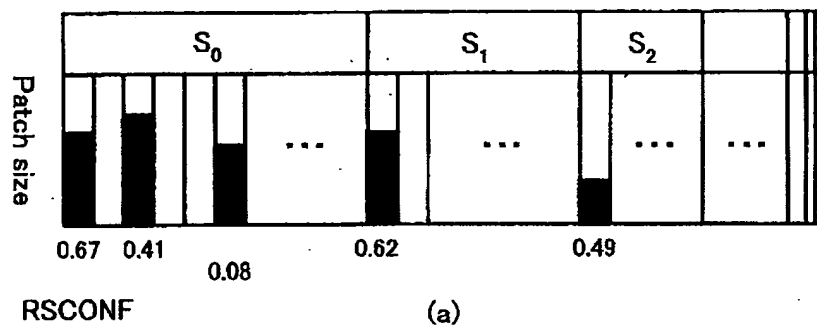


FIG 15

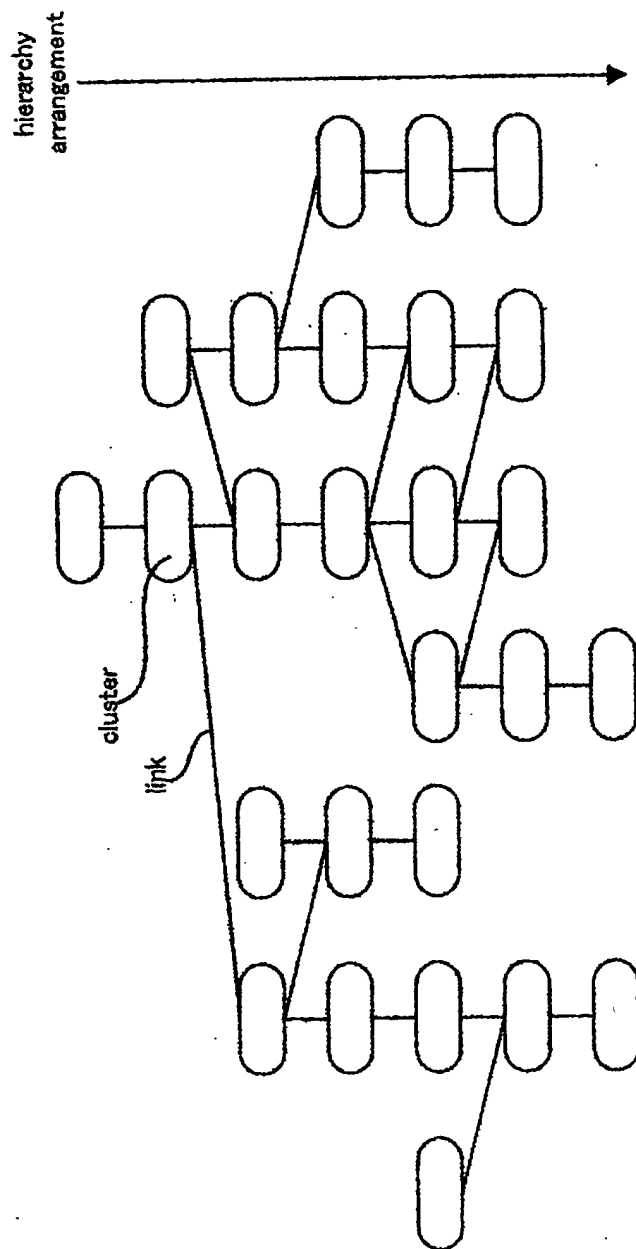


FIG 16

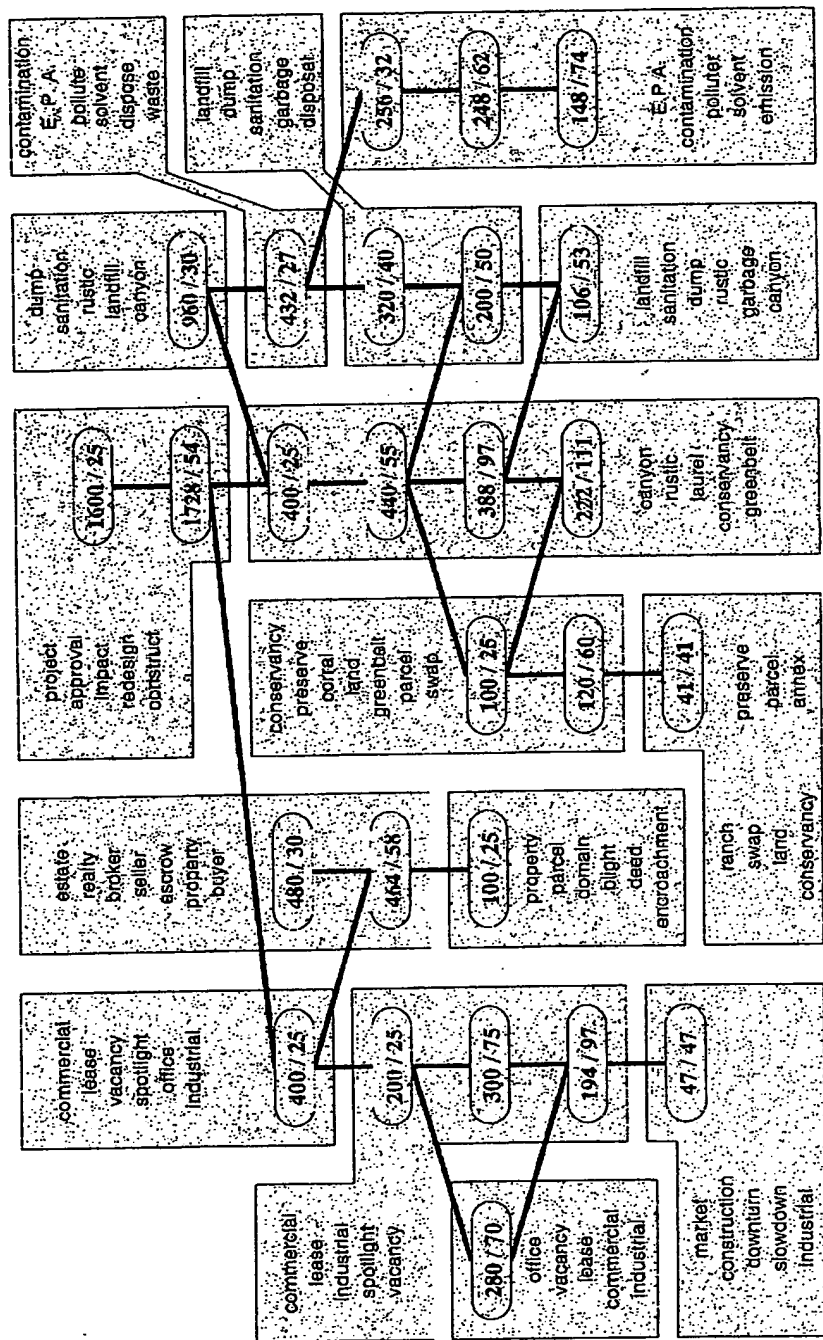


FIG 17

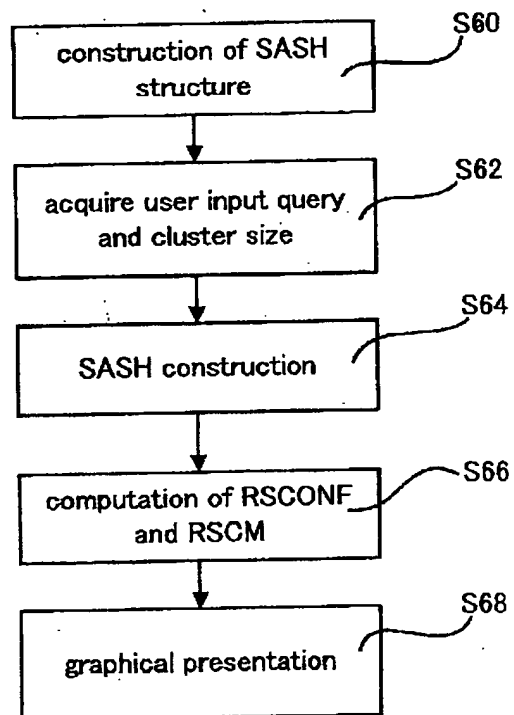


FIG 18

Patches including queried nodes

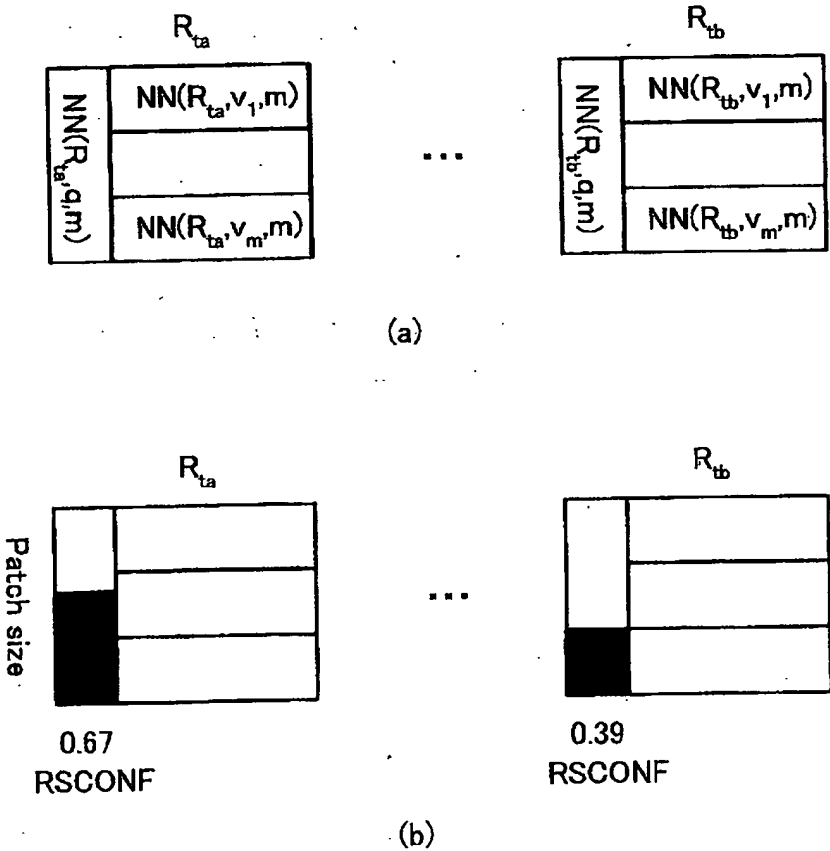


FIG 19

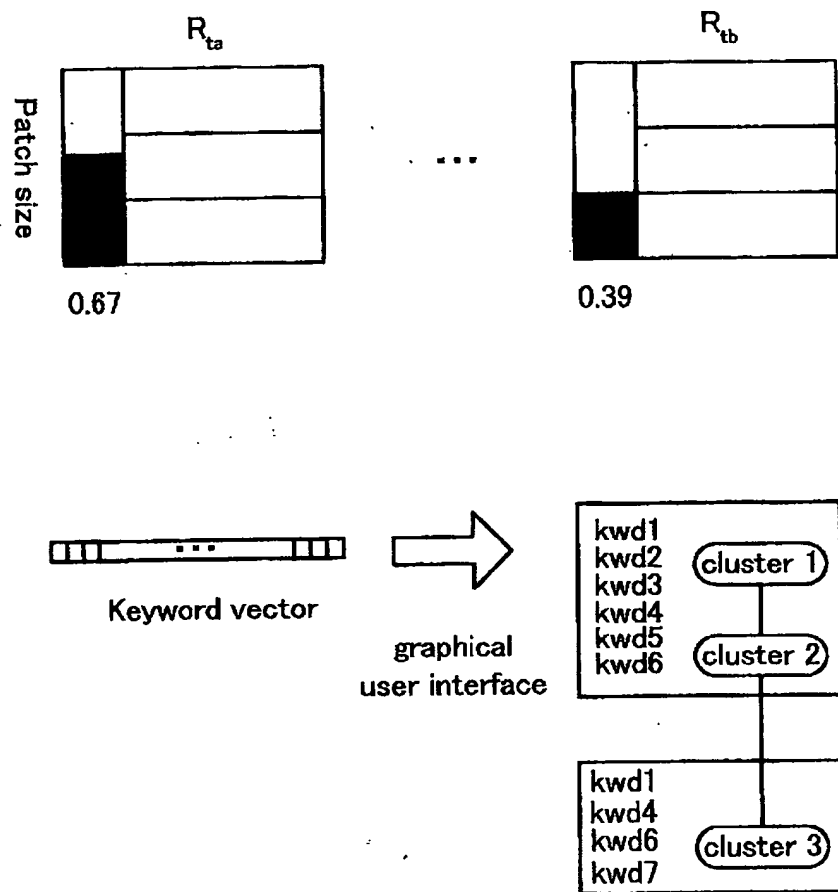


FIG 20

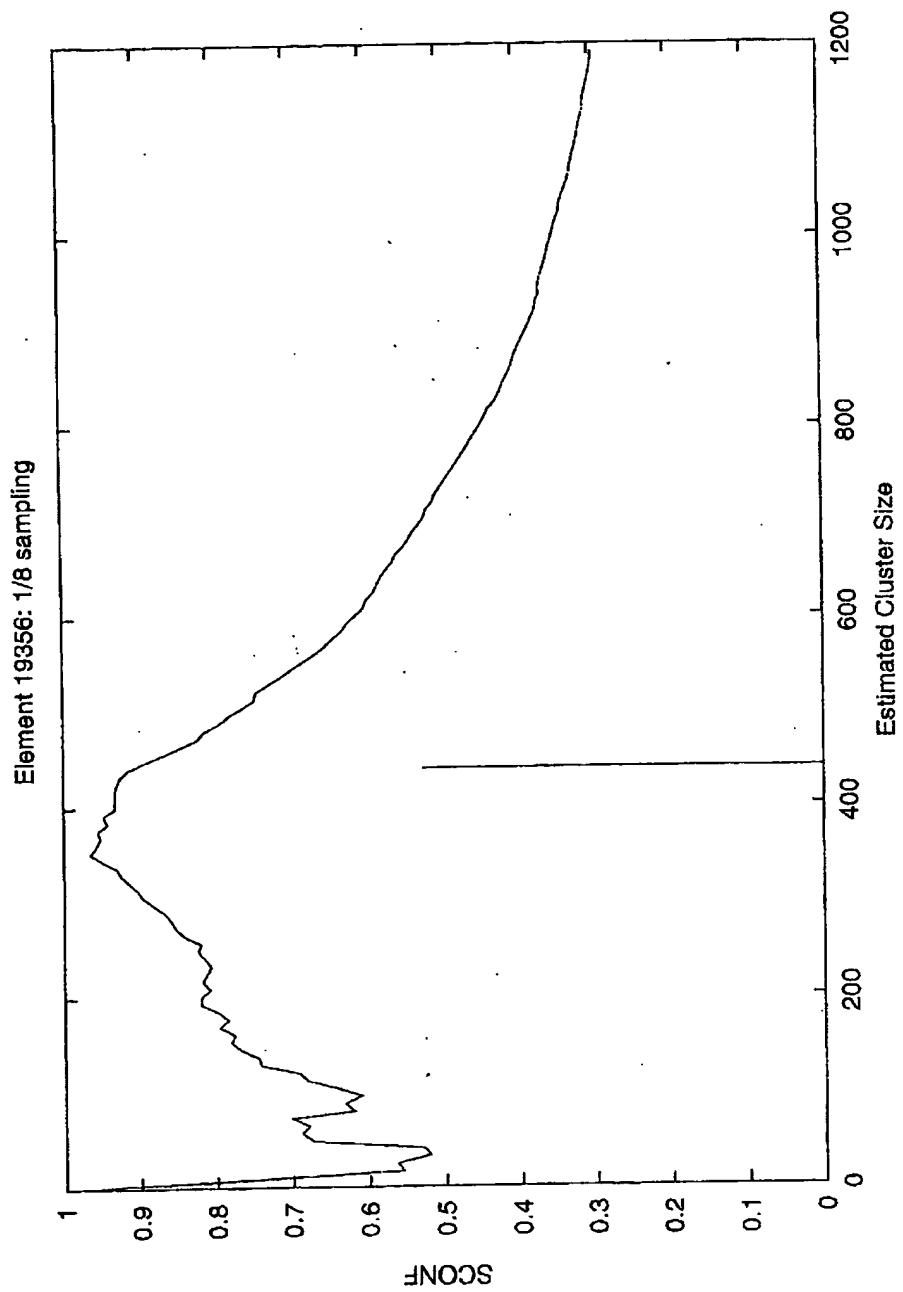


FIG 21